

Playing with the Bandwidth Conservation Law

Orange Labs

Farid Benbadis, Nidhi Hegde, Fabien Mathieu, Diego Perino, Research &
Development
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Bandwidth conservation law in a P2P system?

- Common assumption: download speed = average upload
- And in real life?
 - Joost offers 700kbps download for 120kbps upload (6/1)
 - BitTorrent: 300/1 for some personal experiments...

Sources	Clients	Reçoit à	Envoie à	Restant
106 (712)	8 (171)	131,5 ko/s	789 o/s	1h 24m
375 (4007)	124 (2384)	922,5 ko/s	2,9 ko/s	5m 12s

- But the bandwidth conservation is a ``physical" law
) maybe applying it correctly can give some insight

Roadmap

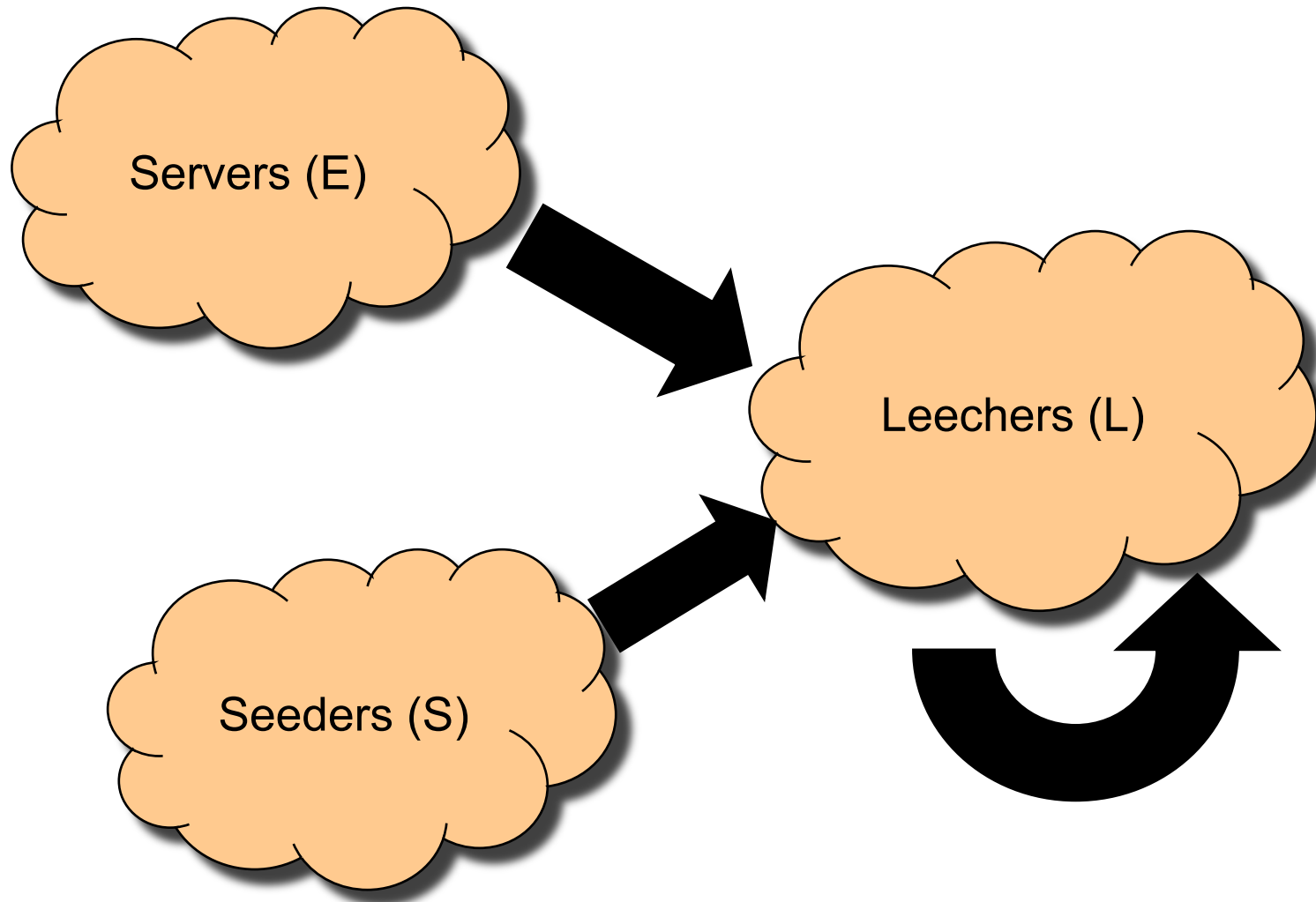
- Model
 - Bandwidth, just bandwidth, only bandwidth
 - Leechers, seeders, servers
- The closed model
 - Some nice formulas
 - Application: is Joost scalable?
- The open model
 - Other formulas (nice too!)
 - Application: maximum download speed in BitTorrent

What matters in a P2P system?

- Access bandwidth
- Core network
- Latency
- Jitter
- Availability
- Allocation schemes
- Routing
- ...

BANDWIDTH

The Leechers, the Seeders and the Servers



Closed Systems, looking for scalability



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A first bandwidth conservation law

- We assume populations and upload capacity are fixed

$$N_L, \bar{u}_L, N_S, \bar{u}_S, U_E$$

- If bandwidth is fairly distributed (same d for all), we have

$$d = \min \left(d_{\max}, \bar{u}_L + \beta \bar{u}_S + \frac{U_E}{N_L} \right),$$

$$\text{with } \beta = \frac{N_S}{N_L}$$

Streaming point of view

- For a given r let $\alpha_X = \frac{\bar{u}_X}{r}$ and $N_E = \frac{U_E}{r}$
- If $d \geq r$, the streamrate can be achieved, otherwise...

- Case 1: $\alpha_L + \beta\alpha_S \geq 1$
 - System is scalable (it works independently of N_L),
 - Servers are unnecessary (except for bootstrap/security)

- Cas 2 : $\alpha_L + \beta\alpha_S < 1$
 - N_L is bounded by $\frac{N_E}{1 - (\alpha_L + \beta\alpha_S)}$
 - ! the server capacity is leveraged by $\frac{1}{1 - (\alpha_L + \beta\alpha_S)}$

Back to Joost

- The upload/download ratio gives $\alpha = \frac{1}{6}$
- If $\beta = 0$ (no seeders), P2P increases the server capacity by 20%
 $\beta = 1$
- If $\beta = 2$, this becomes 50% (most probable behavior)
- If $\beta = 3$, 100%
- ...
- Scalability is achieved only if $\beta \geq 5$ all the time!
- Open question: is Joost more CDN or P2P?

Open Systems

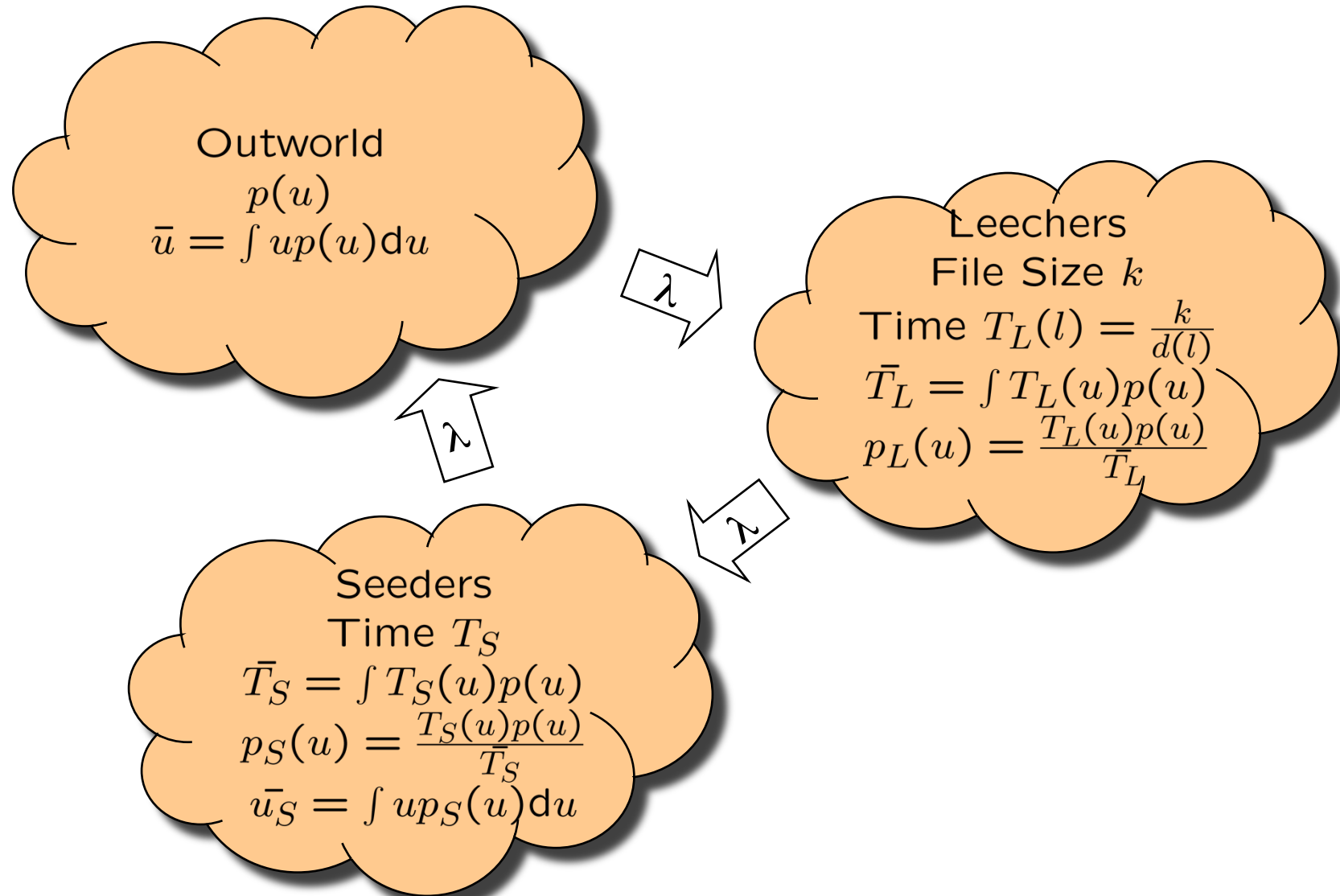
Stationary process



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Leechers/Seeders Life Cycle



Overprovisioning condition

- We use $d = \min \left(d_{\max}, \bar{u}_L + \beta \bar{u}_S + \frac{U_E}{N_L} \right)$, and introduce the intensity.
- This gives a sufficient condition for achieving d_{\max} in any optimal allocation scheme:

$$\bar{T}_S \bar{u}_S + \frac{U_E}{\lambda} > k \left(1 - \frac{\bar{u}}{d_{\max}} \right).$$

- In particular:
 - if $\bar{T}_S \bar{u}_S + \frac{U_E}{\lambda} \geq k$, any target rate can be achieved,
 - if $\bar{T}_S \bar{u}_S \geq k$, any target rate can be achieved, for any λ, U_E .

Back to BitTorrent

- No server and scalability mean that we want the $\bar{T}_S \bar{u}_S \geq k$ condition;
- \rightarrow seeding time must be greater than $\frac{k}{u_S}$;
- In other word, if you seed the time needed to get the file when there is no seeder...
- ...then you download in no time!

Back to BitTorrent

- Interpretation: a system where you pay your bandwidth debt after your download can work.
- Question: how can you force the users to seed?
- First answer: ask politely (sometimes it works)
- Second answer: private tracker and share-ratio policy.

Share-ratio enforcement

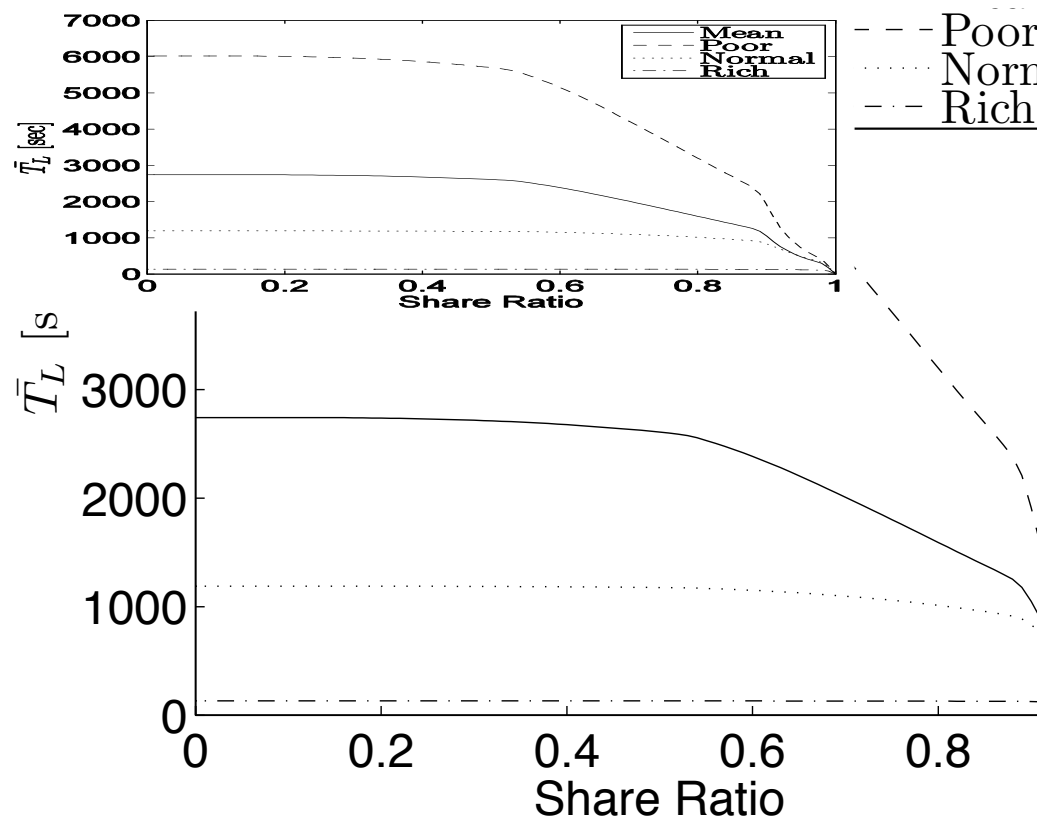
- Some systems force a minimal download/upload share ratio $\tau < 1$ from their participants
- Users that do not meet the share-ratio have "trouble"
- This gives an equation for T_S :

$$T_S(u) = \max\left(0, \frac{\tau k}{u} - T_L(u)\right)$$

- No easy way to solve this in the general case, but one can
 - Consider specific cases,
 - Find singularities,
 - Approximate the solution,
 - Use a numeric solver

) performance studies can be made

Example with heterogeneous bandwidth distribution and Tit-for-Tat



Back to the ``personal experiment''

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- It was a private tracker!
- It also was very popular content...
...but that's another story!

Conclusion

- Bandwidth conservation is simple,
- It applies to all bandwidth-consuming application,
- It gives answers on what we can expect from a P2P (or hybrid) system

- For more applications, play with the article
 - Jetlaging with the bandwidth conservation law,
 - Free-riding with the bandwidth conservation law
 - ...

Thank you very many!



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Bonus 1: worldwide convolution

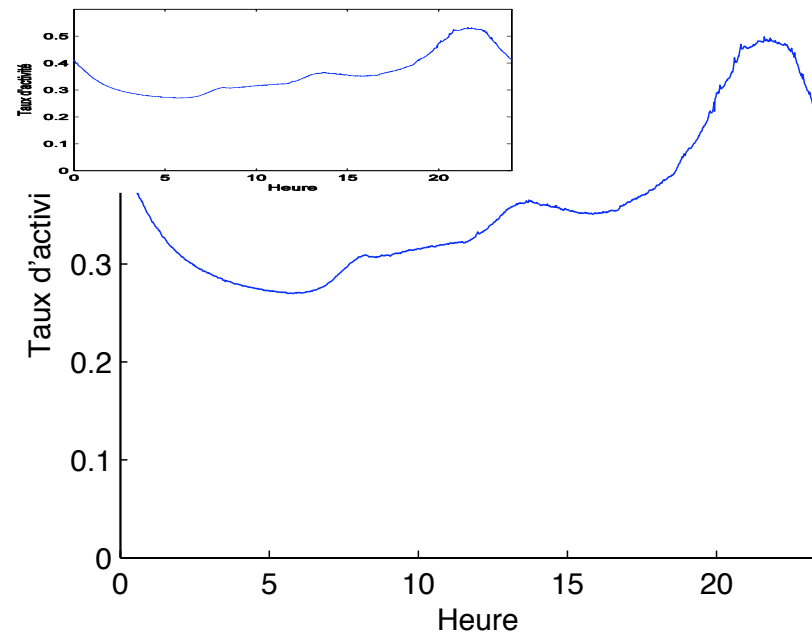


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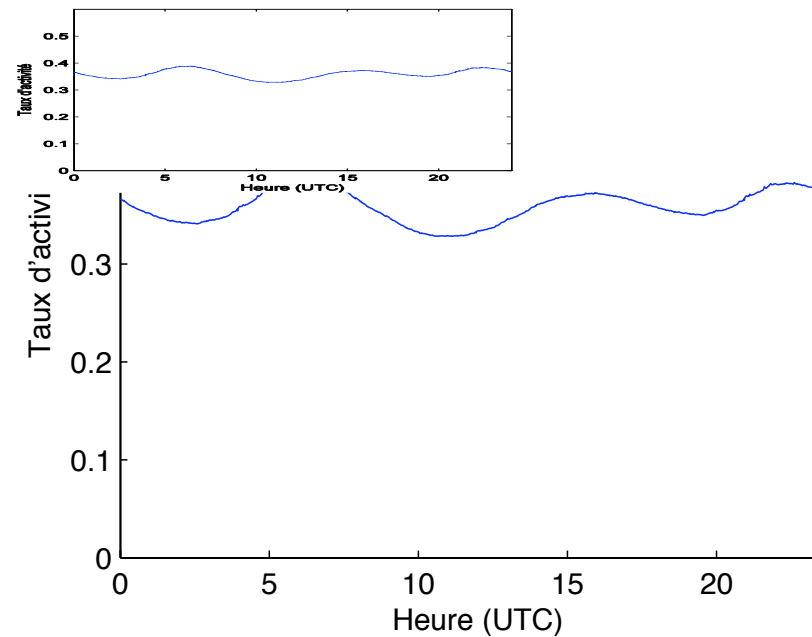
Rush hour

- In practice, leeching (and therefore β) is a function of time
- The dimensioning must be performed in the rush hour



Rush hour – geographical smoothing

- In a worldwide context, the rush hour is timeshifted
- ! Worldwide peaks are smoothed by convolution



Pros and cons of worldwide convolution

- Use the smoothed peak in place of the real one
 - More clients,
 - More streamrate,
 - Less servers
- Each time the local dimensioning rules are broken, extra-bandwidth MUST come from transatlantic/transcontinental links
 - Latency issues,
 - Core networks gets loaded,
 - Isn't managing local connections the current network trend?

Bonus 2: streaming in the open model



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Target streamrate

- We want to know if a given rate r can be achieved. . .
- If $\bar{T}_S \bar{u}_S + \frac{U_E}{\lambda} \geq k$, overprovisioning is achieved.
- Otherwise, the achievability condition is

$$\lambda k \leq \frac{U_E}{1 - \alpha_L - \beta \alpha_S}$$

- The intensity $\frac{U_E}{k}$ that can be supported by servers alone is leveraged by $\frac{1}{1 - \alpha_L - \beta \alpha_S}$

Introducing Tit-for-Tat

- For a proper understanding of share-ratio in BitTorrent, we need a Tit-for-Tat model first
- Assume some exchanges are reciprocity-driven
- d is not constant anymore, T_L neither, and $p_L \neq p$
- Several models try to describe Tit-for-Tat. The simplest gives

$$d(u) = \gamma u + (1 - \gamma) \bar{u}_L + \frac{\bar{T}_S}{\bar{T}_L} \bar{u}_S + \frac{U_E}{\lambda \bar{T}_L}$$

- γ is the Tit-for-Tat coefficient
- No easy way to solve this in the general case, but one can
 - Consider specific cases,
 - Find singularities,
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Bonus 3: Freeriders in BitTorrent



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Theoretical application: Freeriders in the Sky

- Assume a proportion p_f of newcomers are freeriders (no upload at all).
- Question: how can a stationary regime exist in presence of freeriders?
- Answer:

$$p_f \leq (1 - \gamma) + \gamma \left(\frac{\bar{T}_S \bar{u}_S}{k} + \frac{U_E}{\lambda k} \right)$$

- Beyond the critical value, the system cannot treat the freeriders, which accumulate